

Unit - II

BOOLEAN ALGEBRA & LOGIC GATES

BOOLEAN THEOREMS

<i>Identity</i>	Dual
Operations with 0 and 1: 1. $X + 0 = X$ (identity) 3. $X + 1 = 1$ (null element)	2. $X.1 = X$ 4. $X.0 = 0$
Idempotency theorem: 5. $X + X = X$	6. $X.X = X$
Complementarity: 7. $X + X' = 1$	8. $X.X' = 0$
Involution theorem: 9. $(X')' = X$	

BOOLEAN LAWS

<p><i>Identities for multiple variables</i></p>	
<p>Cummutative law: 10. $X + Y = Y + X$</p>	<p>11. $X.Y = Y.X$</p>
<p>Associative law: 12. $(X + Y) + Z = X + (Y + Z)$ $= X + Y + Z$</p>	<p>13. $(XY)Z = X(YZ)$ $= XYZ$</p>
<p>Distributive law: 14. $X(Y + Z) = XY + XZ$</p>	<p>15. $X + (YZ) = (X + Y)(X + Z)$</p>
<p>DeMorgan's theorem: 16. $(X + Y + Z + \dots)' = X'Y'Z' \dots$ or $\{f(X_1, X_2, \dots, X_n, 0, 1, +, \dots)\}$ $= \{f(X_1', X_2', \dots, X_n', 1, 0, \dots, +)\}$</p>	<p>17. $(XYZ \dots)' = X' + Y' + Z' + \dots$</p>

De MORGAN'S THEOREM

- DeMorgan's Theorems describe the equivalence between gates with inverted inputs and gates with inverted outputs.
- Simply put, a NAND gate is equivalent to a Negative-OR gate, and a NOR gate is equivalent to a Negative-AND gate.

De MORGAN'S THEOREM

(i)Statement

The theorem states that the complement of sum of variables is equal to the product of their individual complements.

$$\overline{A + B} = \overline{A} . \overline{B}$$

Proof

<i>A</i>	<i>B</i>	\overline{A}	\overline{B}	$\overline{A + B}$	$\overline{A} . \overline{B}$
<i>0</i>	<i>0</i>	<i>1</i>	<i>1</i>	<i>1</i>	<i>1</i>
<i>0</i>	<i>1</i>	<i>1</i>	<i>0</i>	<i>0</i>	<i>0</i>
<i>1</i>	<i>0</i>	<i>0</i>	<i>1</i>	<i>0</i>	<i>0</i>
<i>1</i>	<i>1</i>	<i>0</i>	<i>0</i>	<i>0</i>	<i>0</i>

(ii) Statement-

The theorem states that the complement of product of variables is equal to the sum of their individual complements.

$$\overline{A \cdot B} = \overline{A} + \overline{B}$$

Proof

<i>A</i>	<i>B</i>	\overline{A}	\overline{B}	$\overline{A \cdot B}$	$\overline{A} + \overline{B}$
<i>0</i>	<i>0</i>	<i>1</i>	<i>1</i>	<i>1</i>	<i>1</i>
<i>0</i>	<i>1</i>	<i>1</i>	<i>0</i>	<i>1</i>	<i>1</i>
<i>1</i>	<i>0</i>	<i>0</i>	<i>1</i>	<i>1</i>	<i>1</i>
<i>1</i>	<i>1</i>	<i>0</i>	<i>0</i>	<i>0</i>	<i>0</i>

Another simple way of remembering the theorem is '**Cut the line and change the sign**'.

De Morgan's law is used to simplify the Boolean expressions in digital circuits. De Morgan's laws can be applied to any number of variables.

E.g. De Morgan's laws for three variables

$$\overline{A + B + C} = \overline{A} \cdot \overline{B} \cdot \overline{C} \quad \&$$

$$\overline{A \cdot B \cdot C} = \overline{A} + \overline{B} + \overline{C}$$

Perfect induction

- **perfect induction**, or the **brute force method**, is a method of mathematical proof in which the statement to be proved is split into a finite number of cases or sets of equivalent cases and each type of case is checked to see if the proposition in question holds.
- This is a method of direct proof. A proof by exhaustion contains two stages:
- A proof that the cases are exhaustive; i.e., that each instance of the statement to be proved matches the conditions of (at least) one of the cases.
- A proof of each of the cases.

Induction proof of $x+x'y=x+y$

Use perfect induction to prove $x+x'y=x+y$

x	y	$x'y$	$x+x'y$	$x+y$
0	0	0	0	0
0	1	1	1	1
1	0	0	1	1
1	1	0	1	1

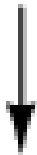
equivalent

Reduction of logical expression using Boolean Algebra

- Boolean algebra finds its most practical use in the simplification of logic circuits.
- If we translate a logic circuit's function into symbolic (Boolean) form, and apply certain algebraic rules to the resulting equation to reduce the number of terms and/or arithmetic operations, the simplified equation may be translated back into circuit form for a logic circuit performing the same function with fewer components.
- If equivalent function may be achieved with fewer components, the result will be increased reliability and decreased cost of manufacture.

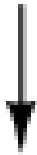
$$AB=1$$

$$A + AB$$



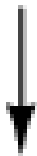
Factoring **A** out of both terms

$$A(1 + B)$$



Applying identity **A + 1 = 1**

$$A(1)$$

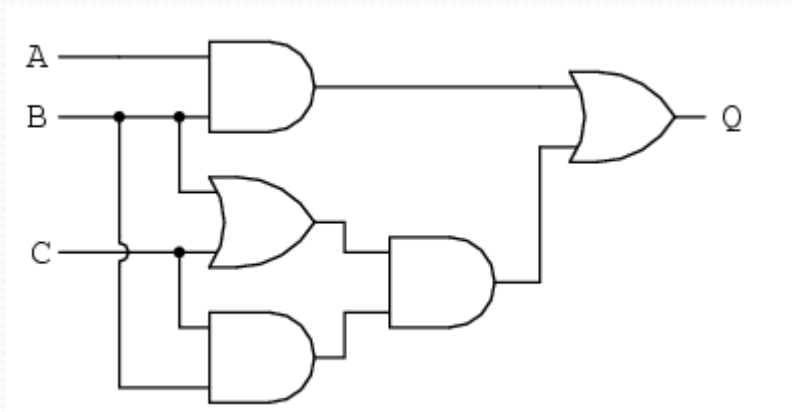


Applying identity **1A = A**

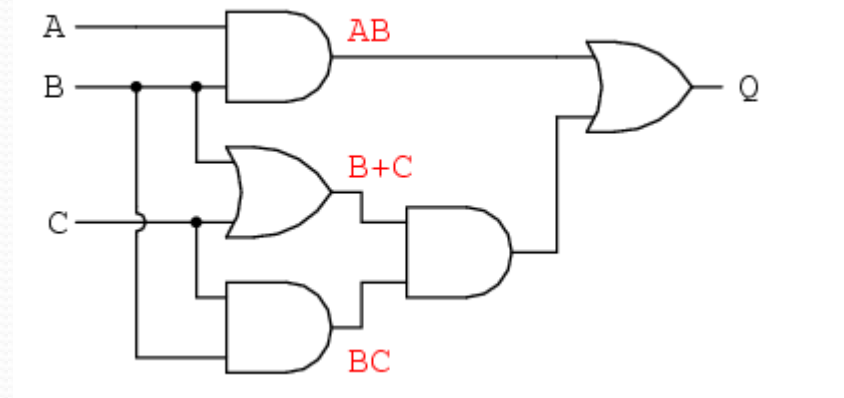
$$A$$

Derive Boolean Expression from circuit

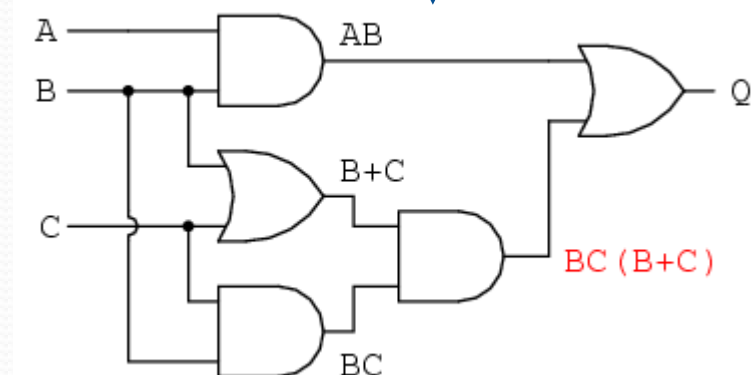
- To understand how to Derive Boolean Expression from circuit, let us consider following circuit.



Step -1



Step - 2



Step -3

